

Quantum Garbled Circuits

Zvika Brakerski

Weizmann

Henry Yuen

Columbia

arxiv.org/abs/2006.01085



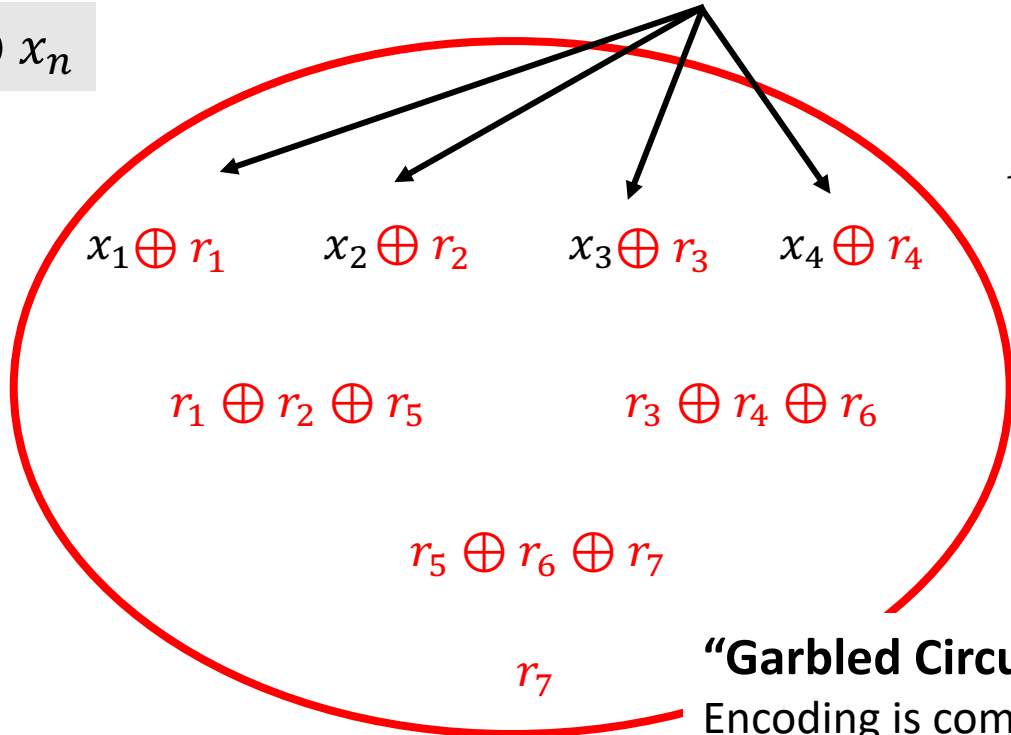
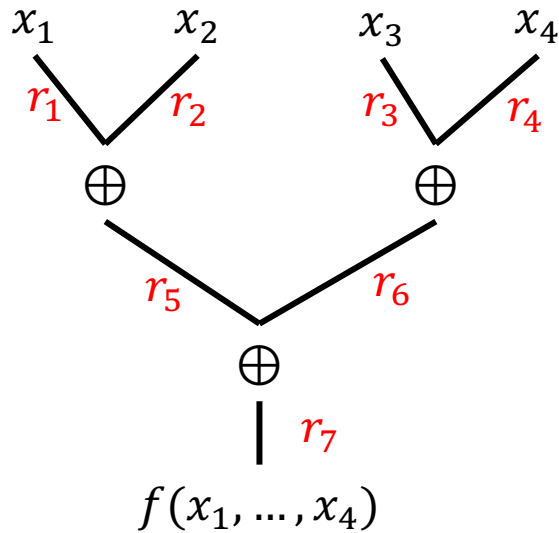
European Research Council
Established by the European Commission

Randomized Encoding (RE) & Garbled decomposable

each input bit treated separately + offline part

Example: Parity $f(x_1, \dots, x_n) = x_1 \oplus x_2 \oplus \dots \oplus x_n$

depth $\log n$



$\hat{f}(x_1, \dots, x_4, r_1, \dots, r_7)$

multiple output bits

depth 2

“Garbled Circuit”

Encoding is computed “gate-by-gate”

NAND gate more complicated

[Yao86, BMR90]

\hat{f} is a **Randomized Encoding (RE)** of f [AIK06]:

- Correctness (decodability): $\hat{f}(x_1, \dots, x_4, r_1, \dots, r_7) \Rightarrow f(x_1, \dots, x_4)$
- Privacy (simulation): $f(x_1, \dots, x_4) \Rightarrow \text{Sample}(\hat{f}(x_1, \dots, x_4, U, \dots, U))$
(or stat/comp indistinguishable)

Why?

One of the most useful notions in crypto
 The “essence” of f in low complexity
 securely compute $\hat{f} \Rightarrow$ securely compute f
 Delegation, non-BB constructions, LBs ...

Our Results

\hat{F} is a QRE of F : (considering auxiliary input)

- Correctness (decodability): $\hat{F}(x, r, e) \Rightarrow F(x)$

- Privacy (simulation): $F(x) \Rightarrow \text{Sample}(\hat{F}(x, U, e))$ (perfect/stat/comp)

- Defining Quantum RE

- Construction of quantum garbled circuits (decomposable QRE):

Perfect Security

Computational Security

Assumption

-
(same)

OWF
(same)

(vs. Classical GC
[BMR90])

RE Complexity

QNC_f^0 , size $\text{poly}(s, 2^{2^d})$

QNC_f^0 , size $\text{poly}(s, \lambda)$

QCircuit w/ size s
& depth d

(decoding is linear)

NC^0 , size $\text{poly}(s, 2^d)$

NC^0 , size $\text{poly}(s, \lambda)$

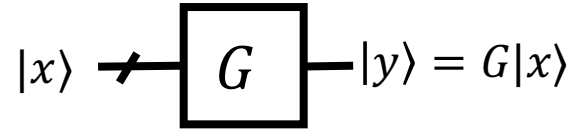
classical input bit \Rightarrow encoded classically

- **Alt. construction:** No shallow encoding, but overall structure is simpler (used in followup [BCKM20]).

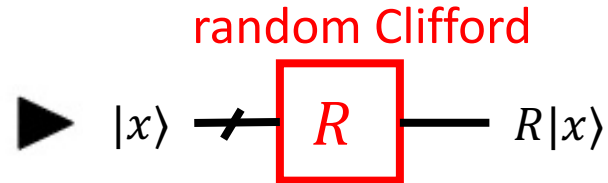
- **Application:** ZK " Σ -Protocol" for QMA w/ favorable properties (comp. to [BG20]).

Warmup: The Group Randomizing QRE

Warmup: Clifford Circuit

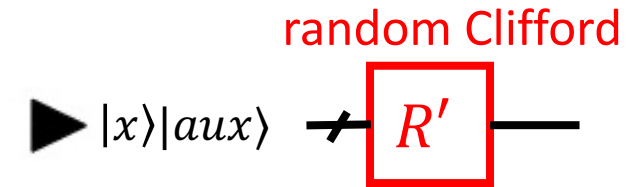


$G \in \text{Clifford}$



$S = GR^{-1}$
canonical circuit
(classical description)

The Goal (\forall Q circuit):

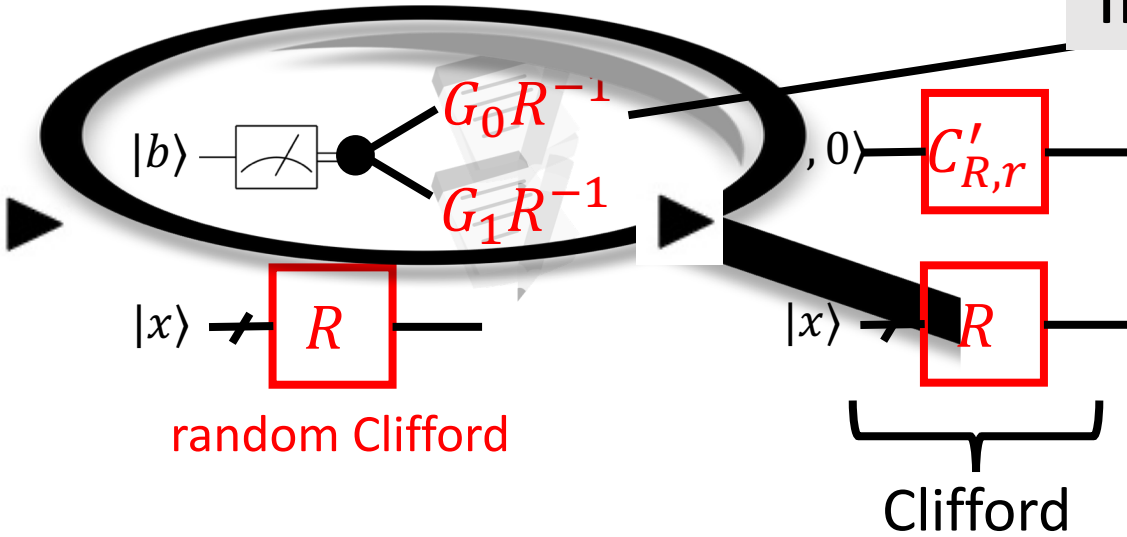
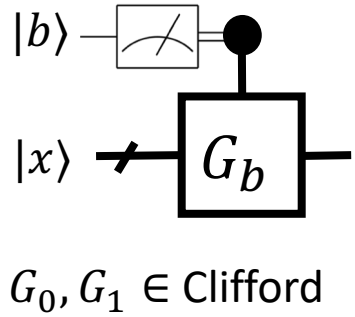


classical side
information

Important feature: Hides G

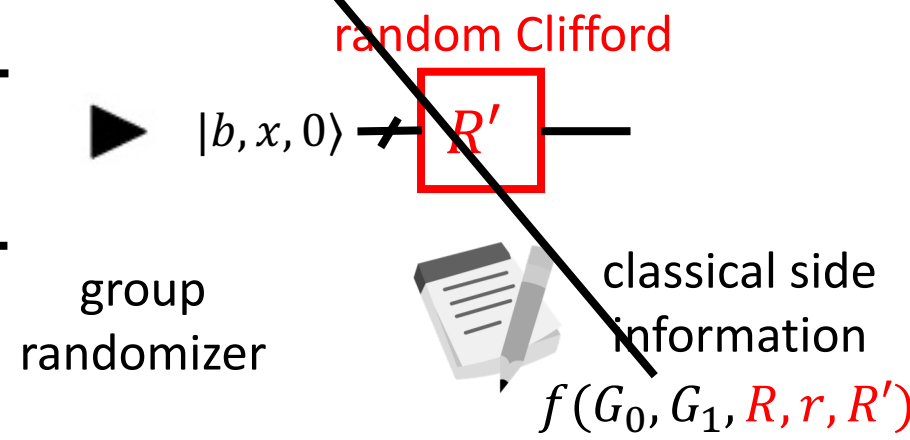
This Work: Quantum RE (QRE)

The C+M Encoder:

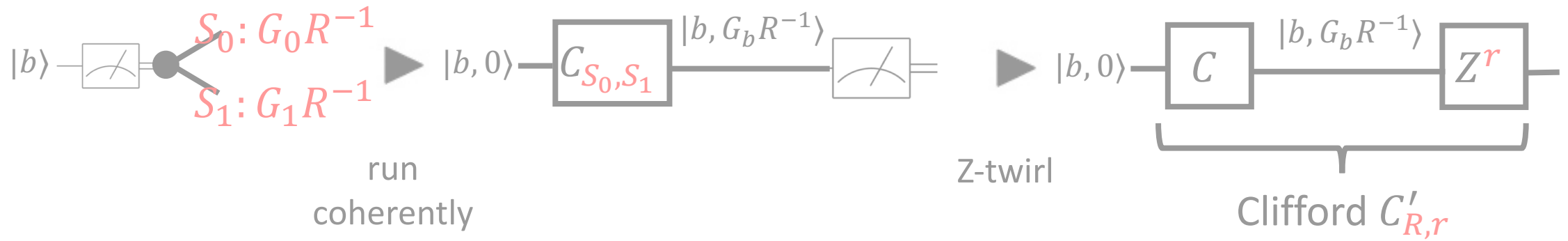


The Goal (\forall Q circuit):

If too complicated, use classical RE to simplify!



Apply recursively
 \Rightarrow QRE of this form \forall Q circuit via [BK05]



Other Applications

Potential applications are numerous:

- Import classical
- New quantum applications?

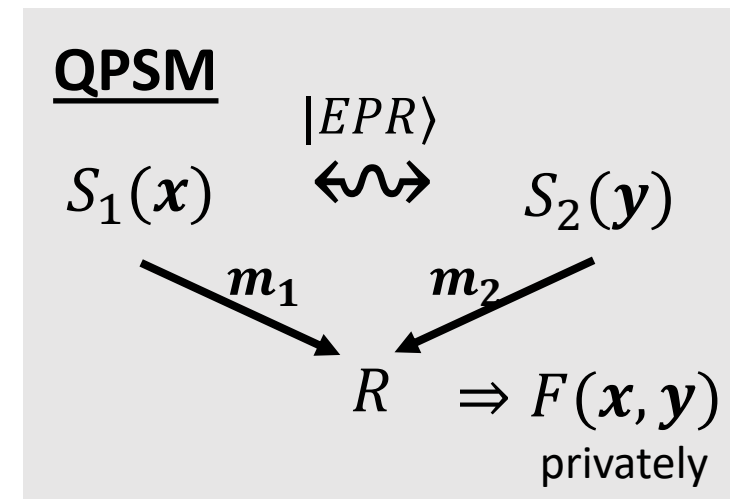
We are unaware of a prior solution, even arbitrarily inefficient

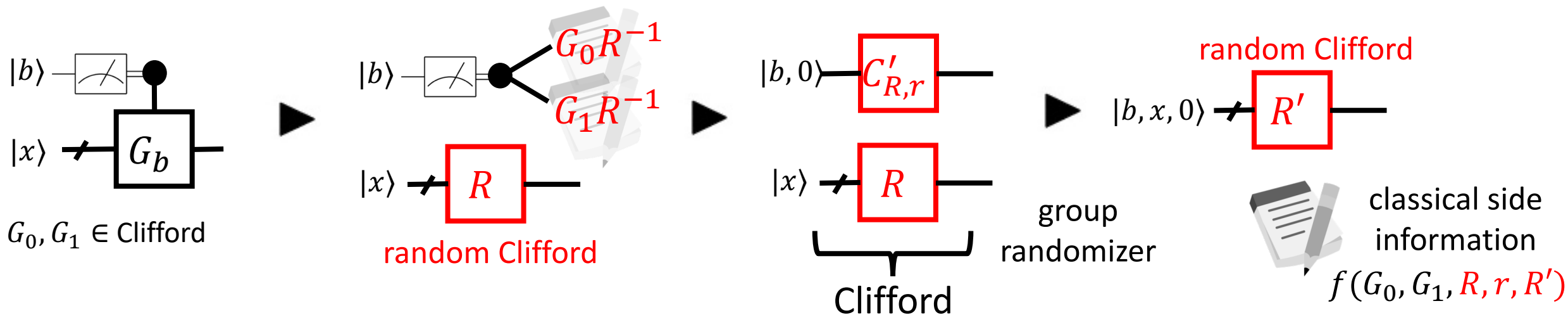
Constant round qMPC?

- Quantum PSM follows by definition.
- Semi-honest 2PC seems straightforward a la Yao.
- Malicious 2PC in 3 rounds recently by [BCKM20].

Functional Encryption and Obfuscation?

- Single-key FE seems straightforward a la [SS10] (but need definition first!).
- Classical RE for Q circuits + Classical obfuscation => Q-Obfuscation.
Beware of barriers, e.g. [Morimae20].





Thank you

arxiv.org/abs/2006.01085

