

Topological obstructions to implementing controlled unknown unitaries

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Outline

Introduction to the problem

Definitions

Task

Our model of computation

Main Result ← **Impossibility of controlled unknown unitary**

Proof of the topological lemma ← **Basic algebraic topology**

Insights ← **An interesting difference between two models**

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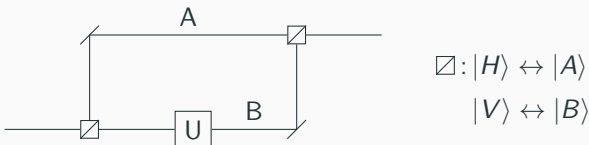
The problem

Given $U \in U(d)$ as a black-box implement

$$\text{control}(U) = |0\rangle\langle 0| \otimes \mathbb{1} + |1\rangle\langle 1| \otimes U$$

Background

- Implementation using Mach-Zehnder interferometer



- ✓ implements $|H\rangle\langle H| \otimes \mathbb{1} + |V\rangle\langle V| \otimes U$
- ✗ **not black-box!** - we know the oracle's location

The problem

Given $U \in U(d)$ as a black-box implement

$$\text{control}(U) = |0\rangle\langle 0| \otimes \mathbb{1} + |1\rangle\langle 1| \otimes U$$

Background

- Implementation using Mach-Zehnder interferometer - **not black box**
- $\text{control}(U)$ impossible [Ara+14; Tho+18]

assume $\text{control}(U)$ possible

U	$\mathbb{1}$	$-\mathbb{1}$
$\text{control}(U)$	$\mathbb{1}_{\text{control}} \otimes \mathbb{1}_{\text{target}}$	$Z_{\text{control}} \otimes \mathbb{1}_{\text{target}}$
$\text{control}(U)(+\rangle \otimes v\rangle)$	$ +\rangle \otimes v\rangle$	$ -\rangle \otimes v\rangle$

but distinguishing $\mathbb{1}$ from $-\mathbb{1}$ is unphysical!

The problem

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Change the task function

$$\text{control} : U(d) \rightarrow L(\mathcal{H}_t) \text{ to } \text{control}_\phi : U(d) \rightarrow L(\mathcal{H}_t)$$

$$\text{where } \text{control}_\phi(U) = |0\rangle\langle 0| \otimes \mathbb{1} + e^{i\phi(U)} |1\rangle\langle 1| \otimes U$$

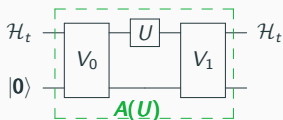
The modified problem

Given $U \in U(d)$ as a black-box implement

$\text{control}_\phi(U) = |0\rangle\langle 0| \otimes \mathbb{1} + e^{i\phi(U)} |1\rangle\langle 1| \otimes U$ for any real function ϕ

Further background

- $\text{control}_\phi(U)$ impossible from one use of U [Ara+14]



$A(U) = V_1 (U \otimes \mathbb{1}) V_0$ linear in U
 \Rightarrow impossibility by linearity arguments

\rightarrow multiple queries to U ?

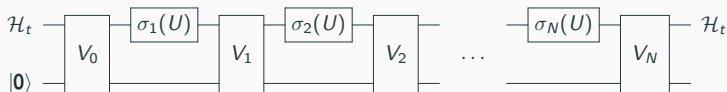
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Further background

- $\text{control}_\phi(U)$ impossible from one use of U [Ara+14]
- process tomography? - **not** control_ϕ
- $\text{control}_\phi(U)$ possible from queries to U^q , $q \in \mathbb{Q}$ [Don+19]



$\sigma_i \in \{U \mapsto U^q\}_{q \in \mathbb{Q}}$

Our problem

task function $U \mapsto \text{control}_\phi(U)$

query functions $\text{id} : U \mapsto U$, $\text{inv} : U \mapsto U^\dagger$

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Implementing controlled unitary is a "Task"

Two kinds of operator functions

- task function $t : U(d) \rightarrow L(\mathcal{H}_t)$ we want to implement
- query function $\sigma_i : U(d) \rightarrow L(\mathcal{H}_i)$ we are allowed to assume, $\sigma_i \in \Sigma$

Definition

Task is a pair (t, Σ)

t task function

Σ query alphabet

Controlled unitary as a Task

- ? $(U \mapsto \text{control}_\phi(U), \{id, inv\})$
- ✓ $(U \mapsto \text{control}_\phi(U), \{U \mapsto U^q\}_{q \in \mathbb{Q}})$
- ✓ $(U \mapsto \text{control}_\phi(U^d), \{id\})$

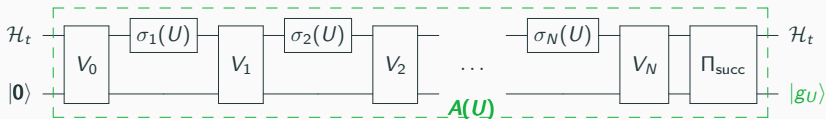
The generalised problem:

- ? Achieve a task $\mathbf{c}\text{-}U^m = (U \mapsto \text{control}_\phi(U^m), \{id, inv\})$ for $m \in \mathbb{Z}$

Our model of computation

Postselection oracle algorithm

- any finite number of queries from Σ
- unitary followed by a binary measurement $\{\Pi_{\text{succ}}, \Pi_{\text{fail}}\}$



- worst-case: $p_{\text{succ}}(U) = \|\lvert g_U \rangle\|^2 > 0$ for all $U \in U(d)$
- represented by $A : U(d) \rightarrow L(\mathcal{H}_t \otimes \mathcal{K})$

ϵ -approximate¹

ϵ -approximates the task (t, Σ) if $A(U) (\mathbb{1} \otimes \lvert 0 \rangle \langle 0 \rvert) \approx_{\epsilon} t(U) \otimes \lvert g_U \rangle \langle 0 \rvert$
for some nonzero $\lvert g_U \rangle$

¹ ϵ measured by diamond norm if $p_{\text{succ}}(U) = 1$, for $p_{\text{succ}}(U) \neq 1$ see arXiv:2011.08487

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The dichotomy

$$c-U^m = (U \mapsto \text{control}_\phi(U^m), \{id, inv\})$$

Theorem

Let $d \in \mathbb{N}$ be the oracle dimension, $m \in \mathbb{Z}$, $\epsilon < \frac{1}{2}$

✓ if $d \mid m$ then $c-U^m$ is

✗ if $d \nmid m$ then $c-U^m$ is not ϵ -approximable.

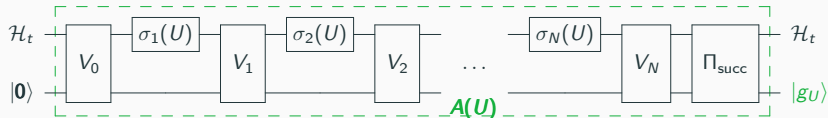
The Topological Lemma

Let $d \in \mathbb{N}$, $m \in \mathbb{Z}$. A continuous m -homogeneous function $U(d) \rightarrow S^1$ exists iff $d \mid m$.

$$f(\lambda U) = \lambda^m f(U)$$

proof of Theorem given The Topological Lemma

$$\mathbf{A} : U(d) \rightarrow L(\mathcal{H}_t \otimes \mathcal{K}) \quad \mathbf{A}(U) = \Pi_{\text{succ}} V_N (\sigma_N(U) \otimes \mathbb{1}) \dots V_1 (\sigma_1(U) \otimes \mathbb{1}) V_0$$



The dichotomy

$$c-U^m = (U \mapsto \text{control}_\phi(U^m), \{id, inv\})$$

Theorem

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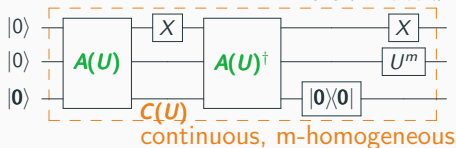
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proof of Theorem given The Topological Lemma

$A : U(d) \rightarrow L(\mathcal{H}_t \otimes \mathcal{K})$ • continuous, w -homogeneous

• $A(U)(\mathbb{1} \otimes |0\rangle\langle 0|) = (|0\rangle\langle 0| \otimes \mathbb{1} + e^{i\phi(U)} |1\rangle\langle 1| \otimes U^m) \otimes |g_U\rangle\langle 0|$



$$f(U) := \frac{\langle 0 | C(U) | 0 \rangle}{\langle 0 | A(U)^\dagger A(U) | 0 \rangle}$$

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The Topological Lemma

Let $d \in \mathbb{N}$, $m \in \mathbb{Z}$. A continuous m -homogeneous function $U(d) \rightarrow S^1$ exists iff $d \mid m$.

The case of even d and odd m (i.e. when the function is odd) can be proven via **Borsuk-Ulam Theorem**

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✗ $control_{\phi}(U)$ for all $U \in U(d)$ impossible even with postselection and **even approximately!** - in particular **via process tomography**:

✓ process tomography approximates $control_{\phi_x}(U)$, $x \sim X$
 where X set of "success" measurement outcomes, $|X| \geq 2$

✗ cannot drop the x -dependence:

cannot rotate continuously to a canonical phase ϕ
 follows from our topological lemma²

! note that our **Postselection oracle algorithm** has one "success" outcome \implies it cannot approximate $control_{\phi_x}(U)$, $x \sim X$

algorithm with measurement	$control_{\phi}$	$control_{\phi_x}, x \sim X$
$\{\Pi_{\text{succ}}, \Pi_{\text{fail}}\}$	✗	✗
$\{\Pi_1, \Pi_2 \dots \Pi_{ X }, \Pi_{\text{fail}}\}$	✗	✓

²we thank an anonymous referee for noting another proof via $SU(d)/Z(SU(d))$ covering spaces.

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- [Don+19] Qingxiuxiong Dong, Shojun Nakayama, Akihito Soeda, and Mio Muraō. “Controlled quantum operations and combs, and their applications to universal controllization of divisible unitary operations”. In: *arXiv preprint arXiv:1911.01645* (2019).
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