



# Efficient unitary designs with a system-size independent number of non-Clifford gates

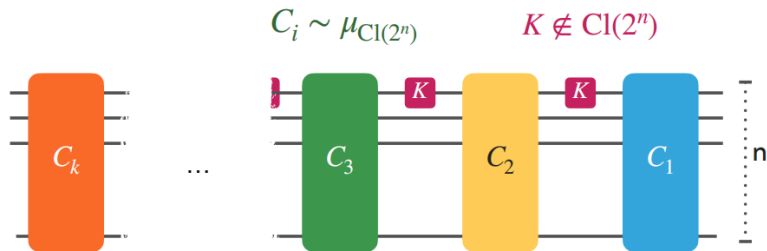
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F. Montealegre-Mora, M. Heinrich, J. Eisert, D. Gross, I. Roth

arXiv:2002.09524

# "Quantum homeopathy": the result in a nutshell



is an **approximate unitary  $t$ -design** in depth  $k \geq O(t^4)$  provided that  $n \geq O(t^2)$ .

# Random unitaries are everywhere

- ▶ **Quantum system identification**: randomized benchmarking, tomography, shadow tomography...
- ▶ Sending information through quantum channels.
- ▶ Models of **quantum information scrambling for black holes**.
- ▶ Generic features of **quantum many-body systems**.

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But: full Haar-randomness requires exponential resources!

# Unitary designs

## Definition

Consider the **moment operator**:

$$\Delta_{\nu,t}(A) := \mathbb{E}_{U \sim \nu} U^{\otimes t} A (U^\dagger)^{\otimes t} \quad (1)$$

**Exact unitary  $t$ -design** is a probability measure  $\nu$  such that

$$\Delta_{\nu,t} = \Delta_{\mu_{\text{Haar}},t}. \quad (2)$$

# Unitary designs

Measures how **evenly spread** a set of unitaries is. E.g. the time evolution under random Hamiltonians.



$t=1$



$t=2$



$t=3$



$t=4$



$t=5$

**Discrete designs** always exist but are complicated!

# Unitary group designs

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Then  $\Delta_{\nu,t}$  is orthogonal projector on **Comm( $U^{\otimes t}$ ,  $U \in G$ )**.

**Example:**  $\Delta_{\mu_{\text{Haar}},t}$  is projector on

$$\text{Comm}(U^{\otimes t}, U \in U(N)) = \text{span}(r(\pi), \pi \in S_t) \quad (\text{Schur-Weyl})$$

# The Clifford group: many faces

- ▶ Generated by Phase gate, Hadamard and controlled NOT:  
 $\text{Cl}(2^n) := \langle \{S, H, CZ\} \rangle$ .
- ▶ Normalizer of the Pauli group:  
 $\text{Cl}(2^n) := \{U \in U(2^n), U\mathcal{P}U^\dagger \subseteq \mathcal{P}\}$ .
- ▶ **Symplectic group** on vector space over finite field.
- ▶ Analogue of **Gaussian operations** for discrete variables.

# The Clifford group: many applications

- ▶ Quantum error correction.
- ▶ Randomized benchmarking.
- ▶ Simulation of quantum circuits.
- ▶ (Shadow) tomography.
- ▶ ...

# The Clifford group as a unitary design

PHYSICAL REVIEW A **96**, 062336 (2017)

## Multiqubit Clifford groups are unitary 3-designs

Huangjun Zhu\*

THE CLIFFORD GROUP FORMS A UNITARY 3-DESIGN

ZAK WEBB<sup>1</sup>

Quantum Information & Computation, 2016

The Clifford group fails gracefully to be a unitary 4-design

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Another face: Multiqubit **Clifford group** is unique non-trivial **3-group** in  $U(2^n)$ .

Bannai, Navarro, Rizo, Tiep, arXiv:1810:02507 (2018)

Guralnick, Tiep, Representation Theory (2005)

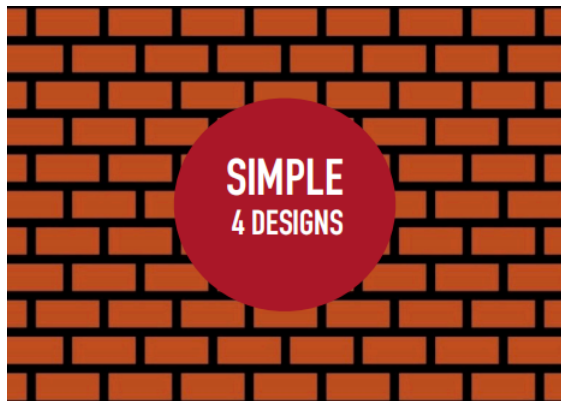
Sawicki, Karnas, Ann. Henri Poincaré (2017).

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$\implies$  4-groups don't exist!

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# Unitary designs via random quantum circuits

## Definition

$\nu$  is  $\varepsilon$ -approximate design if

$$\|\Delta_{\mu_{\text{Haar}},t} - \Delta_{\nu,t}\|_{\diamond} \leq \varepsilon. \quad (3)$$



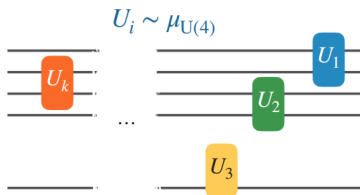
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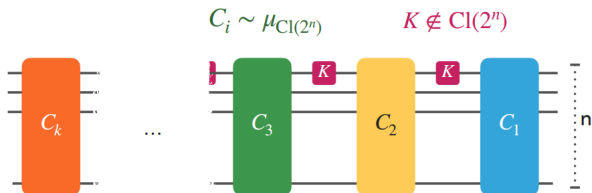
Theorem (Brandão, Harrow, Horodecki)



is an approximate unitary  $t$ -design for  $k \geq O(n^2 t^{9.5})$ .

# Unitary $t$ -designs from random Clifford dominated circuits

## Theorem

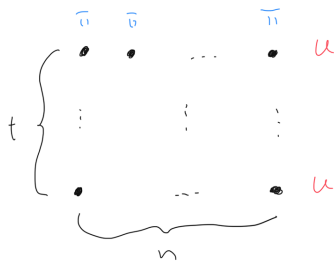


is  $\varepsilon$ -approximate unitary design in depth  $k \geq C_1(K)(t^4 + \log(1/\varepsilon))$   
for all  $n \geq C_2(K)t^2$ .

# Schur-Weyl duality (for the Clifford group)

For  $\pi \in S_t$ , we denote

$$r(\pi) := 2^{-t/2} \sum_{x \in \mathbb{Z}_2^t} |\pi(x)\rangle \langle x|.$$



$$\text{Comm}(U^{\otimes t}) = \text{span}(r(\pi)^{\otimes n})$$

$$\text{Comm}(U^{\otimes t}, U \in \text{Cl}(2^n)) = \text{span}(r(T)^{\otimes n})$$

$T \in \Sigma_{t,t}$  Lagrangian subspaces.

# Proof strategy

Rewrite moment operator

$$\Delta_{\sigma,t} - \Delta_{\mu_{\text{Haar}},t} = \underbrace{[(\Delta_{\mu_{\text{Cl}},t} - \Delta_{\mu_{\text{Haar}}})\text{Ad}_{K^{\otimes t}}]}_{\text{projector}}. \quad (4)$$

- ▶  $\Delta_{\mu_{\text{Cl}},t} - \Delta_{\mu_{\text{Haar}}}$  is projector onto the **orthocomplement** of the **permutations** in the span of **Lagrangian subspaces**.
- ▶ **Technical problem**:  $\{r(T)^{\otimes n}\}$  is not an orthonormal basis.
- ▶ Careful bound on **Gram-Schmidt orthogonalization** of **Lagrangian subspaces** for  $n \geq O(t^2)$ .

# Proof strategy

- ▶ Painful combinatorial argument yields

$$\|\Delta_{\sigma,t} - \Delta_{\mu_{\text{Haar}},t}\|_{\diamond} \leq 2^{O(t^4)+t \log(k)} (1 + 2^{O(t^2)-n})^{5k} \eta^{k-1}.$$

- ▶ Depends on "expectation value" of  $\text{Ad}_{K^{\otimes t}}$  acting on non-permutation Lagrangian subspaces:

$$\eta := \max_{T \in \Sigma_t \setminus S_t} |(r(T) | \text{Ad}_{K^{\otimes t}} | r(T))_{\text{HS}}|.$$

## Lemma

$$\eta \leq 1 - c(K) \log^{-2}(t).$$

# Open problems

- ▶ Can the condition  $n \geq O(t^2)$  be lifted?
- ▶ **Relative** approximate designs. → **partial result**:  $O(n)$  instead of  $O(n^2)$  many non-Clifford gates.
- ▶ Applications of **higher designs**:

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- ▶ Can the condition  $n \geq O(t^2)$  be lifted?
- ▶ **Relative** approximate designs. → **partial result**:  $O(n)$  instead of  $O(n^2)$  many non-Clifford gates.
- ▶ Applications of **higher designs**:
  - ▶ Out-of-time-order correlators and scrambling?
  - ▶ Quantum PUFs?
  - ▶ Equilibration.
  - ▶ Complexity growth.
  - ▶ ...

Mi, et. al., arXiv:2101.08870

Kumar, Mezher, Kashefi, arxiv:2101.05692